The Collo project: An overview



Motivation

 Prototyping new traffic analysis methods is hard!

- System performance strongly depends on traffic patterns
- Long learning phase to fine tune heuristic parameters
- Difficult to properly dimension system without prior knowledge of traffic characteristics



Motivation (cont'd)

- Need for an open network monitoring infrastructure
 - -Widely deployed, diverse datasets
 - -Supports multiple independent users
 - Provides a way to quickly implement the analysis methods
 - Abstracts away the internals of the network monitor



CoMo design goals

- Support multiple arbitrary traffic queries on the packet stream
- Arbitrary means that queries:
 - perform any computation on the packet stream
 - run on past traffic data as well as real-time
 - are coming from different independent users
- Multiple means that queries:
 - compete for resources on the entire infrastructure
 - may be scheduled to run on several systems at once



Related work

- Gigascope (AT&T)
 - GSQL to describe traffic query and schema. Possible to automatically offload to hardware some functions.
- FLAME (UPenn)
 - Focus on safety and trust of in-kernel modules for network monitoring
- Aurora, Borealis (MIT, Brown University)
 - Handle (distributed) continuous queries on data streams
 - Seven operators and automated load shedding techniques
- Scriptroute (UW)
 - Focus on making active measurement simpler to specify and run safely on a distributed architecture



CoMo System model

- Core processes
 - Handle all data movement (wire → memory → disk)
 - Unified interface to packet streams (abstracting capture devices)
 - Schedule queries and optimize allocation of computing resources
 - Maintain full packet trace for a window of time (e.g., 72 hours)
 - Provide query interface and handle user requests to access results



System model (cont'd)

Queries reside in plug-in modules

- Defined by pair <filter:function>
- Filter executed by core processes
- Function to be applied on the packet stream specified in set of callbacks
 - Callbacks used to access query results too
 - Callbacks are closures (i.e. all state is defined in the call) to allow optimization by core processes
- All state is maintained by core processes
 - Allows to stop/pause/resume modules when processing resources are scarce



System architecture



Writing query modules

- Success of system depends on how simple it is to write modules
 - Kismet module (114 semicolons)
 - Top-N destinations (62 semicolons)
 - Bytes/Packets Counter (53 semicolons)
- Each module maintains its own database
 - Using store(), load(), print() callbacks
 - Storage manager deals with indexing data and providing fast disk access.
 - Database schema known only to the module



Historical queries

- Re-use computations made by modules

 (only alternative is to go to packet trace)
- Modules provide a replay() callback
 - Runs on the module's database and generates a synthetic packet stream
 - Comes with a description of the stream
- Core system provides the best match
 - Find modules with all needed information
 - Choose based on storage size and processing costs of replay()



Historical queries (cont'd)

- Replay() allows significant reduction in disk reads at the cost of CPU cycles
 - Assumption is that CPU is not scarce while bandwidth from disk to memory is
- No explicit knowledge of what the two modules are actually doing!
- This method is used today to run on Dante's NetFlow data.

 No need to modify queries that operate on live packet stream



Hardware offload

- Off-loading some module callbacks to hardware.
- Easy for the filter, hard for callbacks
- Examples
 - load() code sent to active storage systems (e.g., DIAMOND) to search the database when no index exists
 - If update() is simple enough it can be compiled to run directly on an IXP's microengine.



Robustness

- Degrade module performance in presence of traffic anomalies or high query load
 - Module chosen based on resource usage
- Gracefully as module resource usage increases
 - Delay results (modules runs offline)
 - Sample input packet stream
 - Last resort: stop the module and resume it later
- Incentives for users to
 - Provide efficient implementations
 - Ask as few resources as possible
 - Implement replay() callback
 - if computation is reusable, it has higher priority

Achieving Safety

- Module may use disproportionate amount of resources
 - Same problem with legit modules and anomalous incoming traffic
- Module may corrupt other module's data
 - Removing pointer arithmetic would solve the problem but it's not feasible
 - Use approach similar to FLAME (based on Cyclone)
- Module may break usage policy
 - e.g., break anonymization scheme
 - Local anonymization is easy to do but not good enough for most users
 - Distributed anonymization still an open problem...



Achieving Scalability

Running on cluster of nodes
 – "similar" modules run on same node



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Making it distributed...

- Demands for global resource management
- Online queries: Optimal Network-wide Sampling
 - Select the nodes that need to run a specific module to minimize the total number of packet processed
 - Module's utility function that depends on the aggregate sampling rate across the network of monitors
 - Optimal algorithm for a class of utility functions
- Retrospective queries: MIND distributed index
 - Nodes share and distribute some information on the traffic they observe
 - Query first look in the index to find nodes that may contain data of interest
 - Prototype implemented on PlanetLab

Conclusion

- CoMo is a platform for fast prototyping traffic analysis methods
- Implementation publicly available
 - Runs on x86 and StrongARM;
 - Supports many network capture devices (DAG, libpcap, Prism2 wireless cards, NetFlow)
 - Full packet capture on Gigabit links
 - Four code releases so far...

 Looking for users and sites where to develop new nodes

More info

Source code available at http://como.sourceforge.net

como-users@lists.sourceforge.net for comments, questions, etc.

